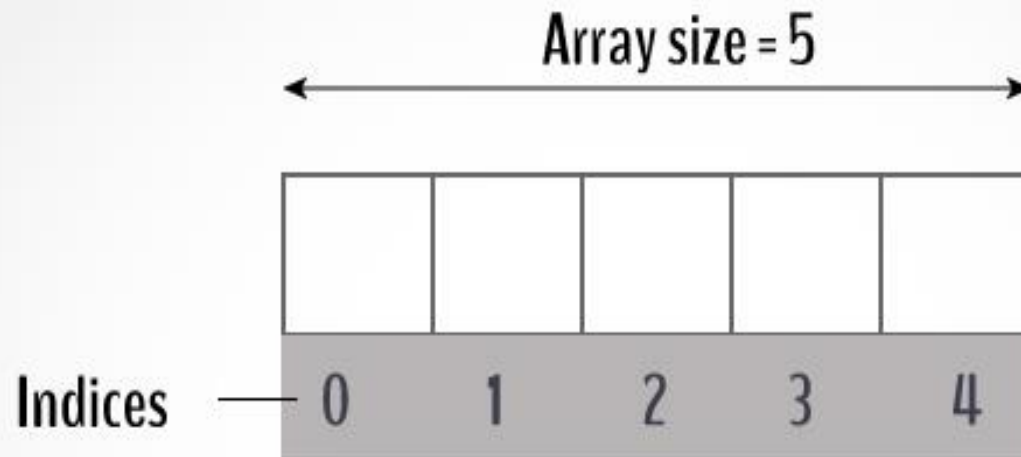


Arrays

- An Array consists of Collection of elements in same data type.
- Elements are stored in consecutive memory locations.
- An array is a set of pairs, index and value. For each index which is defined, there is a value associated with that index.



C Arrays

Array

Structure ARRAY(value, index)

 Declare CREATE()->array

 RETRIEVE(array,index)->value

 STORE(array,index,value)->array;

 for all $A \in \text{array}$, $i, j \in \text{index}$, $x \in \text{value}$ let

 RETRIVE(CREATE,i)::=error

 RETRIVE(STORE(A,i,x),j)::=

 If EQUAL(i,j) then x else RETRIVE(A,J)

 End

End ARRAY

- The function CREATE()
 - It produces a new empty array.



- RETRIVE(array,index)->value
 - It takes input an array and an index, and either returns the appropriate value or an error.

A	10	20	30	40	50
	0	1	2	3	4

- STORE(array,index,value)->array
 - It is used to enter new index pairs.

ORDERED LISTS

- Collection of elements are arranged in sequential order.
- Example
 - Days of the week
 - Values in a card deck
 - $(a_1, a_2, a_3, \dots, a_n)$

Operations

- Find the length of the list, n .
- Read the list from left to right.
- Retrieve the i -th element from the list.
- Store a new value into the i -th position
- Insert a new element at position i .
- Delete the element at position i .

LIST ADT

- A sequence of zero or more elements

$$A_1, A_2, A_3, \dots A_N$$

- N : length of the list
- A_1 : first element
- A_N : last element
- A_i : position i
- If $N=0$, then empty list
- Linearly ordered
 - A_i precedes A_{i+1}
 - A_i follows A_{i-1}

Operations

- printList: print the list
- makeEmpty: create an empty list
- find: locate the position of an object in a list
 - list: 34,12, 52, 16, 12
 - find(52) → 3
- insert: insert an object to a list
 - insert(x,3) → 34, 12, 52, x, 16, 12
- remove: delete an element from the list
 - remove(52) → 34, 12, x, 16, 12
- findKth: retrieve the element at a certain position

Implementation of ADT

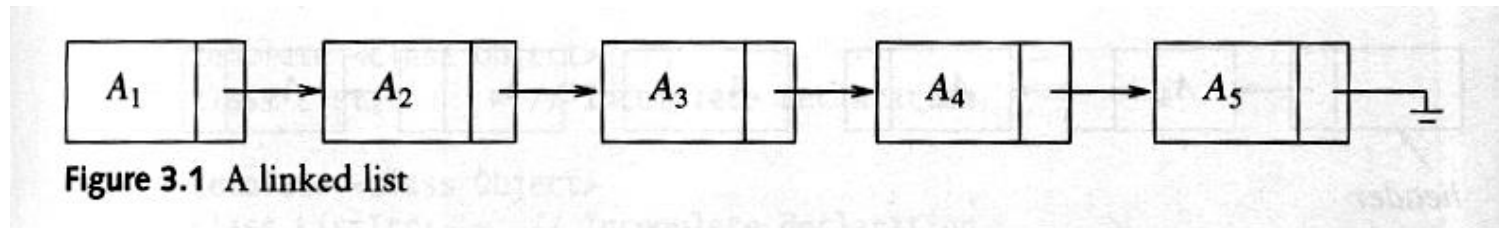
- Choose a **data structure** to represent the ADT
 - E.g. arrays, records, etc.
- Each operation associated with the ADT is implemented by one or more subroutines
- Two standard implementations for the list ADT
 - Array-based
 - Linked list

Array Implementation

- Requires an estimate of the maximum size of the list
 - waste space
- printList and find: linear
- findKth: constant
- insert and delete: slow
 - e.g. insert at position 0 (making a new element)
 - requires first pushing the entire array down one spot to make room
 - e.g. delete at position 0
 - requires shifting all the elements in the list up one
 - On average, half of the lists needs to be moved for either operation

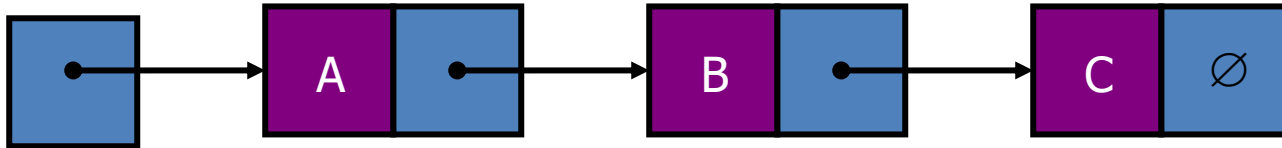
Pointer implementation(Linked List)

- Ensure that the list is not stored contiguously
 - use a linked list
 - a series of structures that are not necessarily adjacent in memory

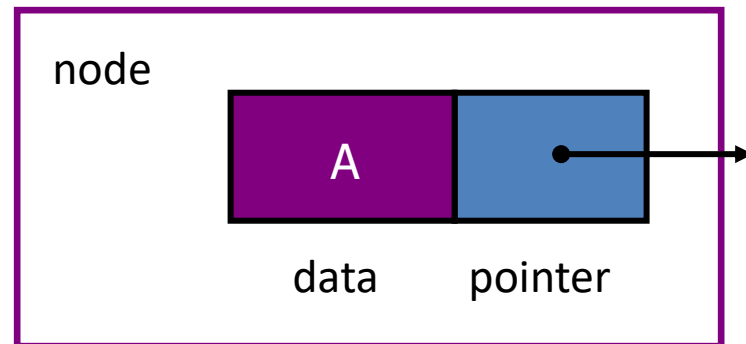


- Each node contains the element and a pointer to a structure containing its successor
 - the last cell's next link points to NULL
- Compared to the array implementation,
 - ✓ the pointer implementation uses only as much space as is needed for the elements currently on the list
 - but requires space for the pointers in each cell

Linked List



- *A linked list* is a series of connected *nodes*
- Each node contains at least
 - A piece of data (any type)
 - Pointer to the next node in the list
- *Head*: pointer to the first node
- The last node points to `NULL`



A Simple Linked List Class

- We use two classes: **Node** and **List**
- Declare `Node` class for the nodes
 - data: `double`-type data in this example
 - next: a pointer to the next node in the list

```
class Node {  
public:  
    double data;           // data  
    Node*      next;       // pointer to next  
};
```

A Simple Linked List Class

- Declare `List`, which contains
 - `head`: a pointer to the first node in the list.
Since the list is empty initially, `head` is set to `NULL`
 - Operations on `List`

```
class List {
public:
    List(void) { head = NULL; }           // constructor
    ~List(void);                          // destructor

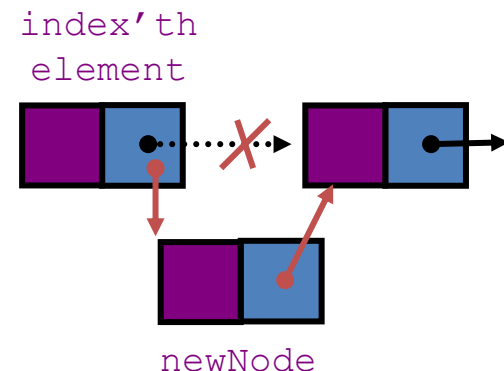
    bool IsEmpty() { return head == NULL; }
    Node* InsertNode(int index, double x);
    int FindNode(double x);
    int DeleteNode(double x);
    void DisplayList(void);
private:
    Node* head;
};
```

Operations of List

- `IsEmpty`: determine whether or not the list is empty
- `InsertNode`: insert a new node at a particular position
- `FindNode`: find a node with a given value
- `DeleteNode`: delete a node with a given value
- `DisplayList`: print all the nodes in the list

Inserting a new node

- `Node* InsertNode(int index, double x)`
 - Insert a node with data equal to `x` after the `index`'th elements. (i.e., when `index = 0`, insert the node as the first element; when `index = 1`, insert the node after the first element, and so on)
 - If the insertion is successful, return the inserted node.
Otherwise, return `NULL`.
(If `index` is `< 0` or `> length` of the list, the insertion will fail.)
- Steps
 1. Locate `index`'th element
 2. Allocate memory for the new node
 3. Point the new node to its successor
 4. Point the new node's predecessor to the new node



Inserting a new node

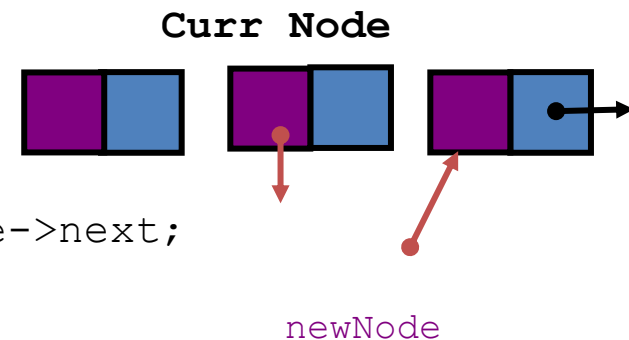
- Possible cases of `InsertNode`
 1. Insert into an empty list
 2. Insert in front
 3. Insert at back
 4. Insert in middle
- But, in fact, only need to handle two cases
 - Insert as the first node (Case 1 and Case 2)
 - Insert in the middle or at the end of the list (Case 3 and Case 4)

Inserting a new node

```
List::InsertNode(int index, double x) {  
    if (index < 0) return NULL;  
  
    int currIndex = 1;  
    Node* currNode = head;  
    while (currNode && index > currIndex) {  
        currNode = currNode->next;  
        currIndex++;  
    }  
    if (index > 0 && currNode == NULL) return NULL;
```

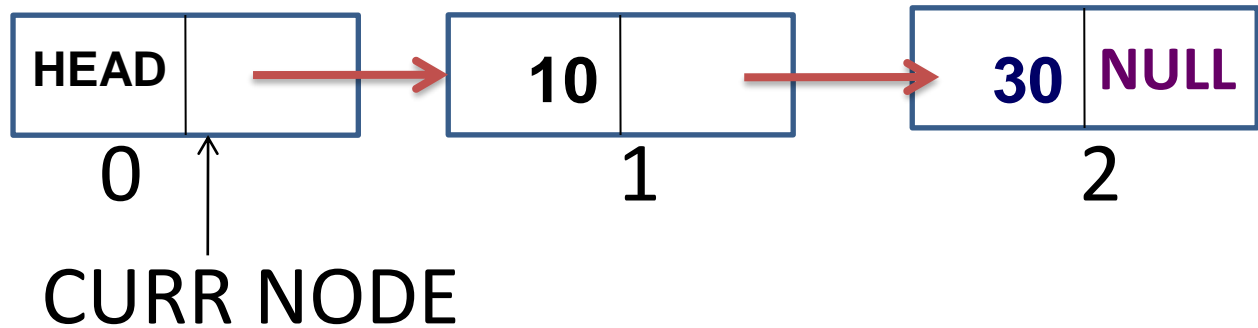
Try to locate index'th node. If it doesn't exist, return NULL.

```
    Node* newNode = new Node;  
    newNode->data = x;  
    if (index == 0) {  
        newNode->next = head;  
        head = newNode;  
    }  
    else {  
        newNode->next = currNode->next;  
        currNode->next = newNode;  
    }  
    return newNode;  
}
```



Inserting a new node

Example: InsertNode(int 1, double 10)



```
Node* newNode = new Node 

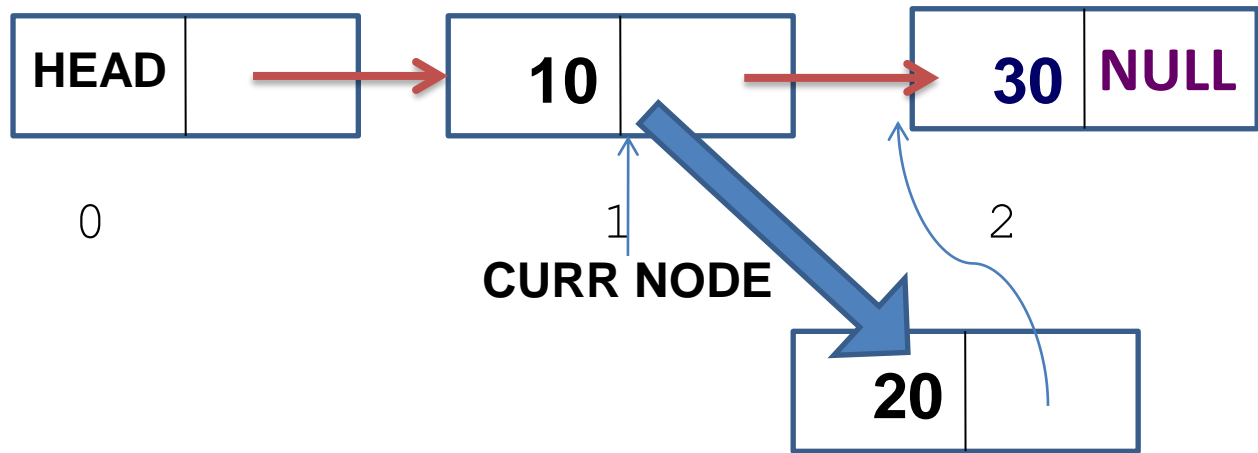
|  |  |
|--|--|
|  |  |
|--|--|

  
newNode->data = 20; 

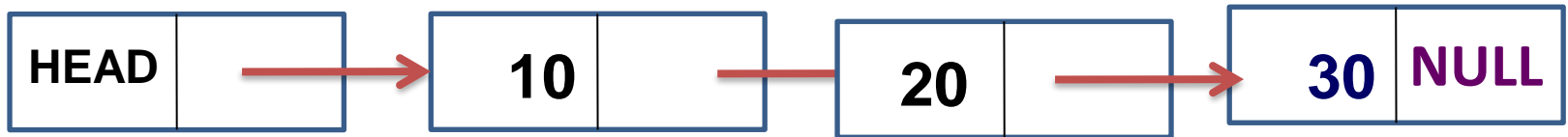
|    |  |
|----|--|
| 20 |  |
|----|--|


```

- `newNode->next = currNode->next;`



- `currNode->next = newNode;`



Finding a node

- `int FindNode(double x)`
 - Search for a node with the value equal to `x` in the list.
 - If such a node is found, return its position. Otherwise, return 0.

```
int List::FindNode(double x) {
    Node* currNode = head;
    int currIndex = 1;
    while (currNode && currNode->data != x) {
        currNode = currNode->next;
        currIndex++;
    }
    if (currNode) return currIndex;
    return 0;
}
```

Deleting a node

- `int DeleteNode(double x)`
 - Delete a node with the value equal to x from the list.
 - If such a node is found, return its position. Otherwise, return 0.
- Steps
 - Find the desirable node (similar to `FindNode`)
 - Release the memory occupied by the found node
 - Set the pointer of the predecessor of the found node to the successor of the found node
- Like `InsertNode`, there are two special cases
 - Delete first node
 - Delete the node in middle or at the end of the list

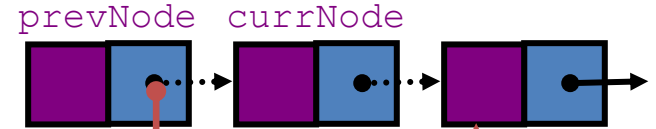
Deleting a node

```
int List::DeleteNode(double x) {
    Node* prevNode = NULL;
    Node* currNode = head;
    int currIndex = 1;
    while (currNode && currNode->data != x) {
        prevNode = currNode;
        currNode = currNode->next;
        currIndex++;
    }
    if (currNode) {
        if (prevNode) {
            prevNode->next = currNode->next;
            delete currNode;
        }
        else {
            head = currNode->next;
            delete currNode;
        }
    }
    return currIndex;
}
return 0;
}
```

Try to find the node with its value equal to x

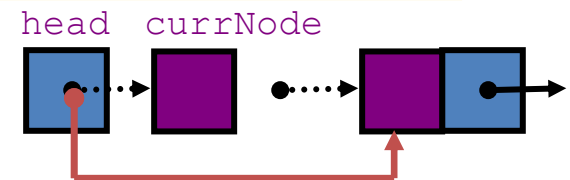
Deleting a node

```
int List::DeleteNode(double x) {
    Node* prevNode = NULL;
    Node* currNode = head;
    int currIndex = 1;
    while (currNode && currNode->data != x) {
        prevNode = currNode;
        currNode = currNode->next;
        currIndex++;
    }
    if (currNode) {
        if (prevNode) {
            prevNode->next = currNode->next;
            delete currNode;
        }
        else {
            head = currNode->next;
            delete currNode;
        }
        return currIndex;
    }
    return 0;
}
```



Deletion

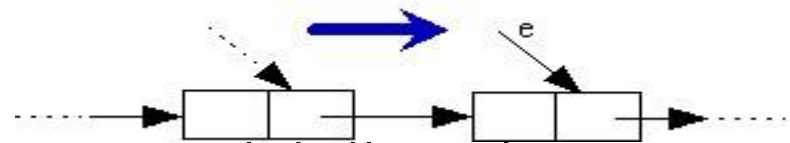
```
int List::DeleteNode(double x) {
    Node* prevNode = NULL;
    Node* currNode = head;
    int currIndex = 1;
    while (currNode && currNode->data != x) {
        prevNode = currNode;
        currNode = currNode->next;
        currIndex++;
    }
    if (currNode) {
        if (prevNode) {
            prevNode->next = currNode->next;
            delete currNode;
        }
        else {
            head = currNode->next;
            delete currNode;
        }
        return currIndex;
    }
    return 0;
}
```



Traversal

Begin at the first node, then follow each *next* reference until the traversal condition is satisfied or until you come to the end.

```
e = e.next;
```



To move an **Element** reference **e** from one node to the next use:

```
public int countNodes() {  
    int count = 0;  
    Element e = head;  
    while(e != null) {  
        count++;  
        e = e.next;  
    }  
    return count;  
}
```

Example: Count the number of nodes in a linked list.

Printing all the elements

- `void DisplayList(void)`
 - Print the data of all the elements
 - Print the number of the nodes

```
void List::DisplayList()
{
    int num          = 0;
    Node* currNode   = head;
    while (currNode != NULL){
        cout << currNode->data << endl;
        currNode     = currNode->next;
        num++;
    }
    cout << "Number of nodes in the list: " << num << endl;
}
```